

# Optimized Distributed Implementation of Multiparty Interactions with Observation

Saddek Bensalem, Marius Bozga, [Jean Quilbeuf](#) and Joseph Sifakis

Verimag

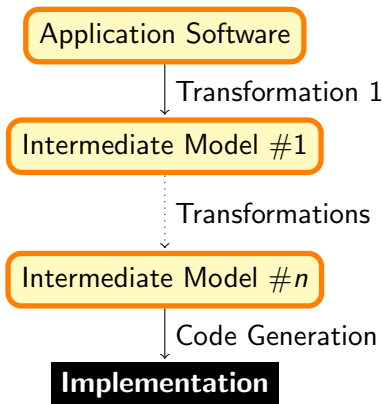
AGERE!12

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Tucson, Arizona, USA

# Rigorous System Design

Sequence of transformations guaranteeing that essential requirements are met by the implementation.



## Model-Based

- Unified semantic model
- Abstraction adapted to task (verification, code generation...)

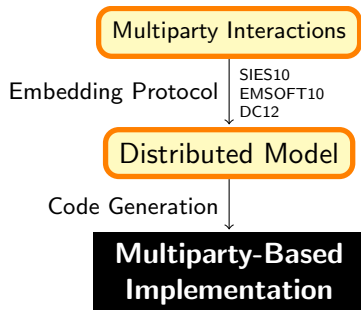
## Component-Based

- Encompass heterogeneity
- Reusability

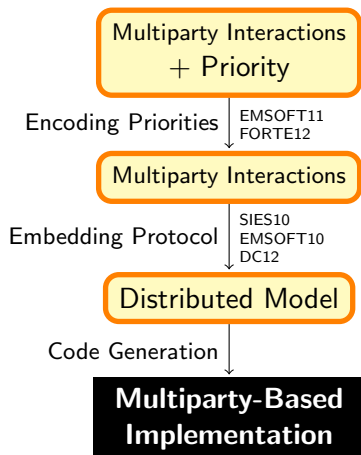
## Correct-by-Construction

- Functional properties preserved by transformations
- Avoid a posteriori verification

# Distributed BIP Design Flow



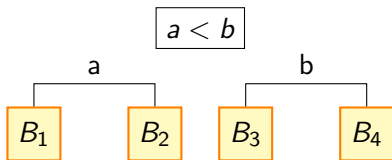
# Distributed BIP Design Flow



# Limits of Multiparty Interactions for encoding Priority

Assume components communicating through Multiparty Interactions.

Example of model with Priority:



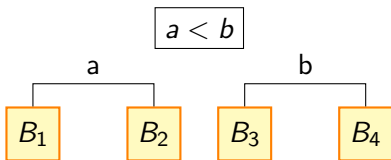
$a$  can execute only if  $b$  cannot.

- Execution of  $a$  does not change the state of  $B_3$  and  $B_4$ .
- Execution of  $a$  depends on the state of  $B_3$  and  $B_4$ .

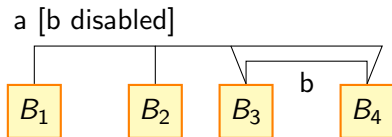
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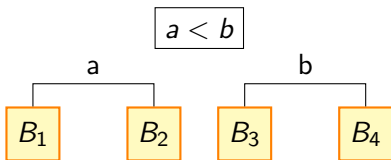
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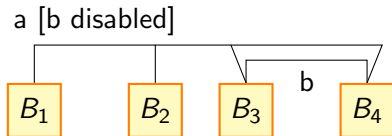
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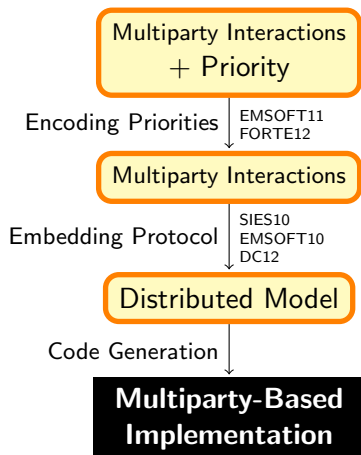
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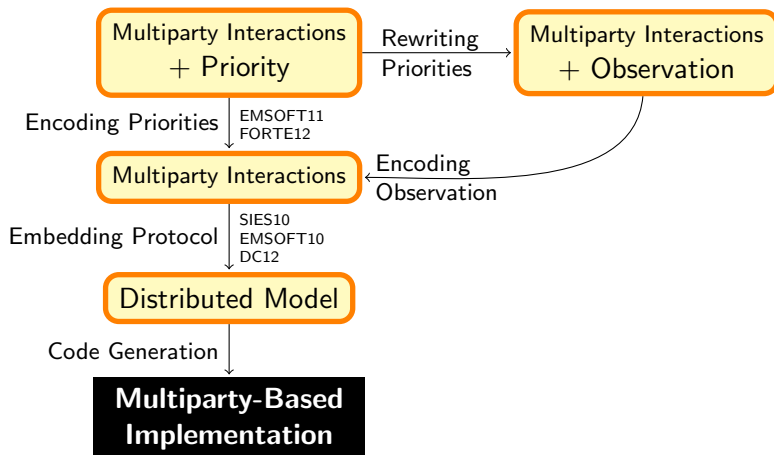
With Multiparty Interaction encoding,  $B_3$  and  $B_4$  become participants in  $a$ .

**Need an intermediate between participation and non-participation.**

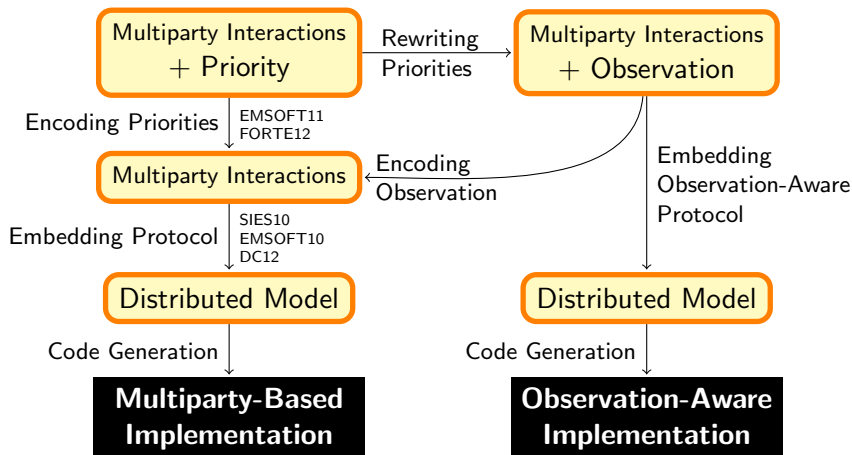
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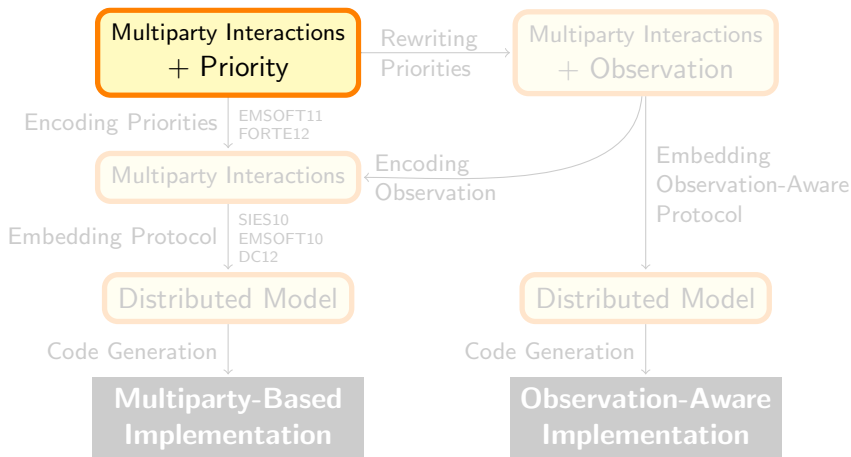
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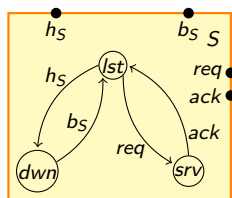
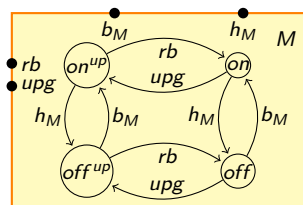


# Outline

- 1 Introduction
- 2 The BIP Model
- 3 Alternative Model: Observation
- 4 Distributed Protocols
- 5 Experiments
- 6 Conclusion

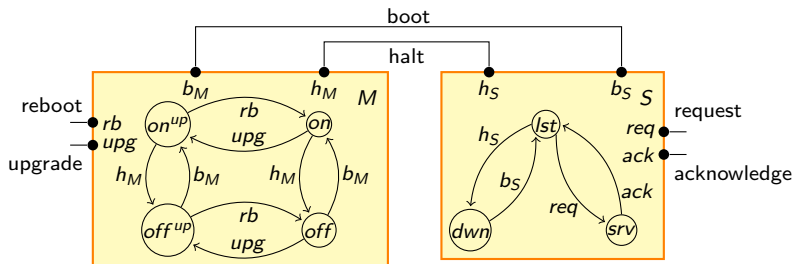
## 2 The BIP Model





## Behavior

- Ports (interface)
- Automaton

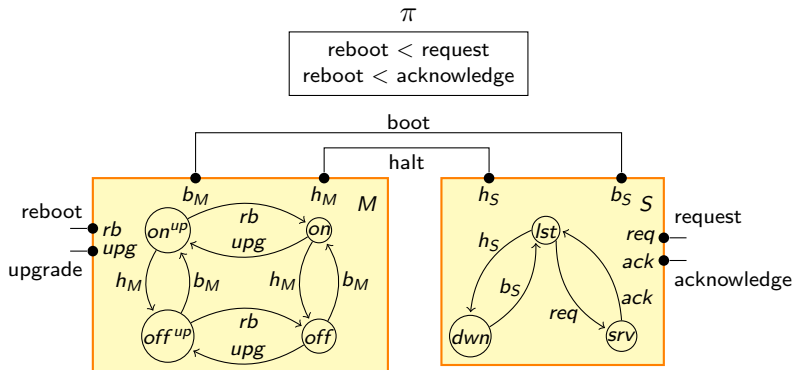


## Behavior

- Ports (interface)
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## Interaction

- Set of ports



## Behavior

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## Priority

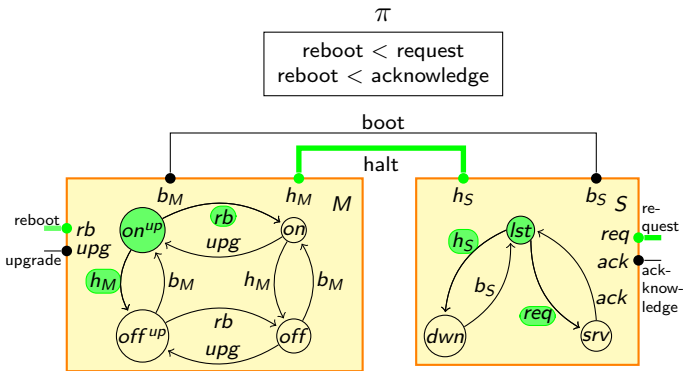
- Partial order over interactions

# BIP Semantics

At global state  $q$ , interaction  $a$  can execute if:

- $a$  is enabled, that is all its ports are enabled, denoted  $EN_a(q)$
- there is no enabled interaction with higher priority

Execution of  $a$  is atomic and changes only the state of the participants.

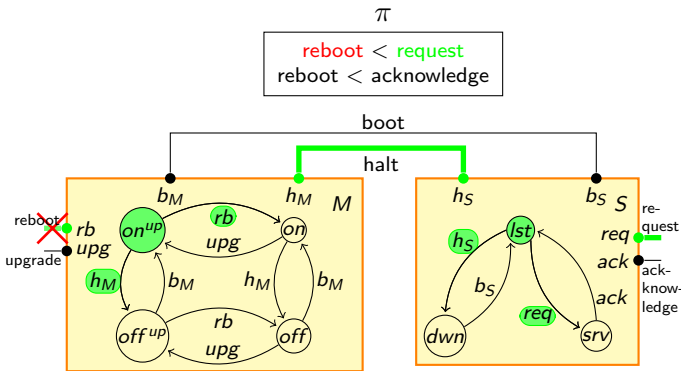


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Enabled interactions at global state  $(on^{up}, lst)$ :

reboot  
request  
halt

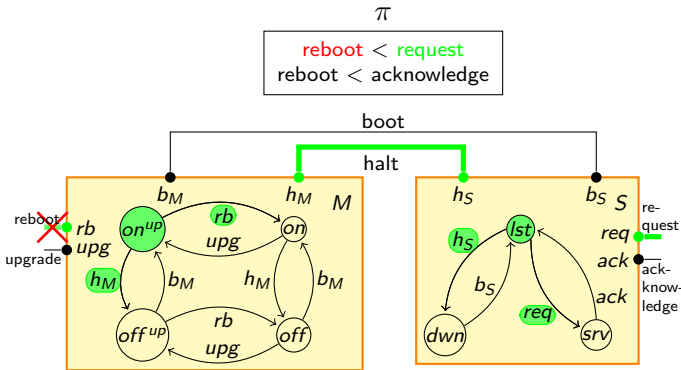
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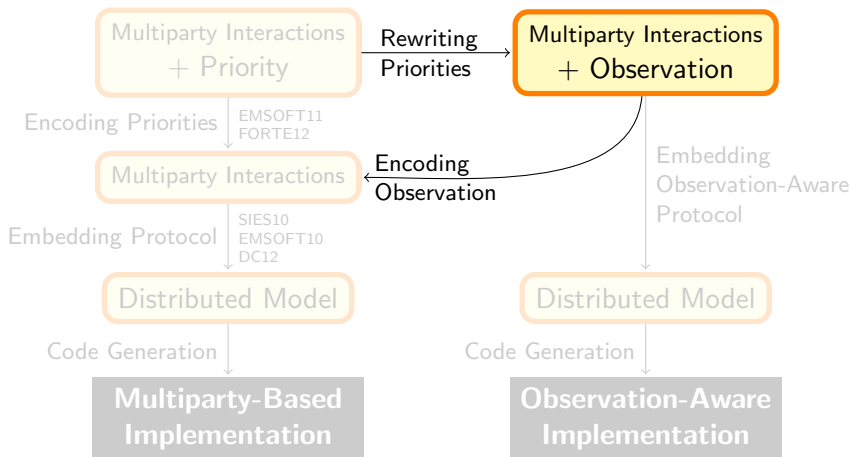
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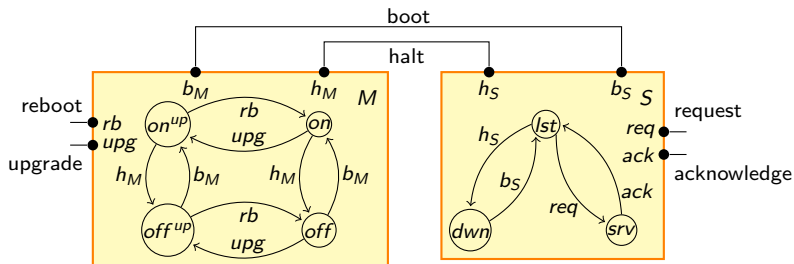
- reboot
- request
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reboot cannot execute because it is dominated by request

Non-deterministic choice between request and halt.

### 3 Alternative Model: Observation

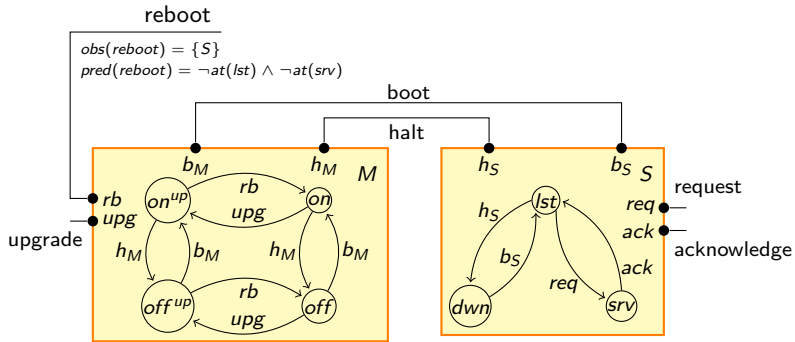




## Behavior

## Interaction

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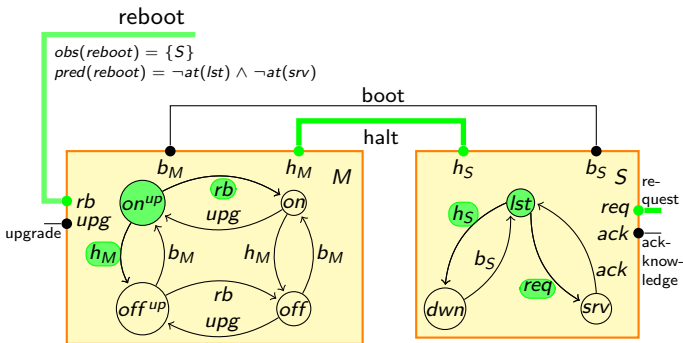
## Observation (each interaction)

- observes a set of components
- is guarded by a state predicate

# BIO Semantics

At global state  $q$ , interaction  $a$  can execute if:

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- the predicate  $pred(a)$  evaluates to *true* at  $q$



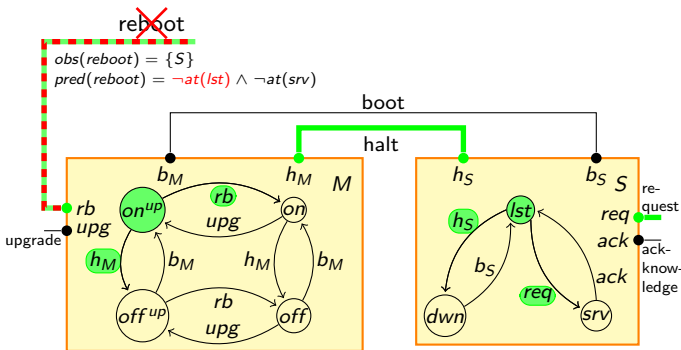
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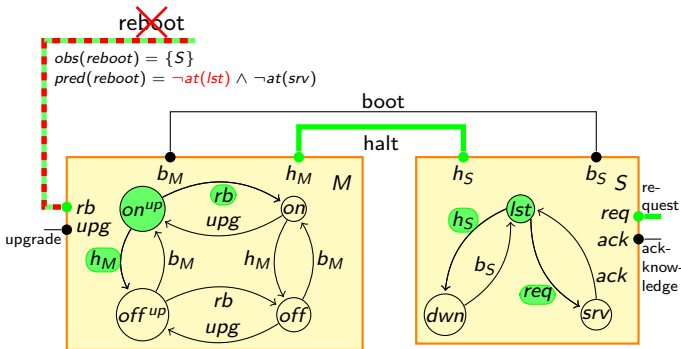
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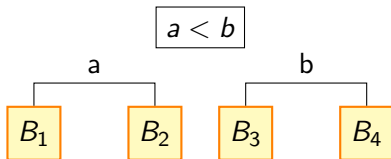
Non-deterministic choice between request and halt.

# Expressing Priority with Observation

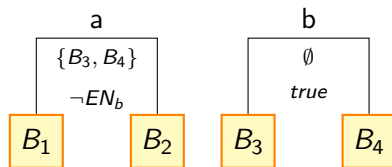
Given a priority  $\pi$ , we define the Observation  $\mathcal{O}_\pi$ , such that each low priority interaction

- observes all components involved in higher priority interactions,
- checks disableness of all higher priority interactions.

**Priority  $\pi$ :**



**Observation  $\mathcal{O}_\pi$ :**

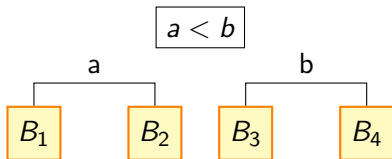


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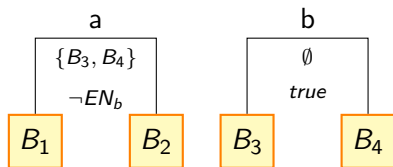
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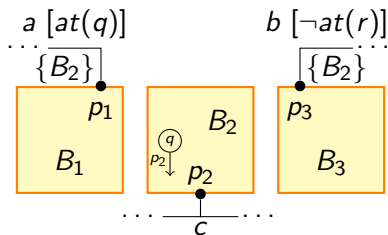
**Proposition** *The two models allow exactly the same transitions (are bisimilar).*

# Expressing Observation with Multiparty Interactions

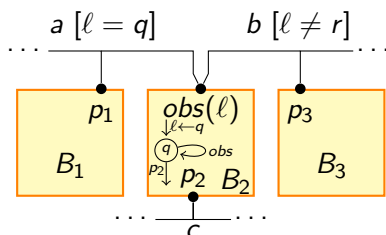
Given model with Observation  $\mathcal{O}$  we define a model with Multiparty Interactions:

- each observed component exports its state through a new port *obs*,
- interactions are extended to port *obs* of the observed component and the predicate is transformed into a guard.

**Observation  $\mathcal{O}$ :**



**Multiparty Interactions:**

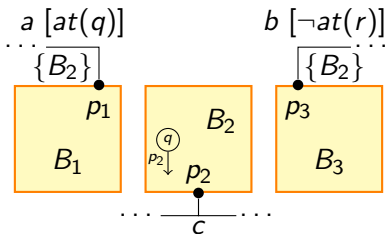


# Expressing Observation with Multiparty Interactions

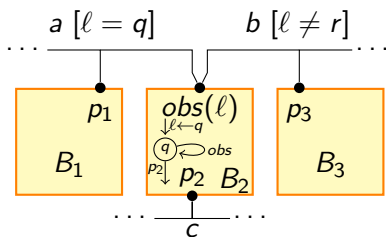
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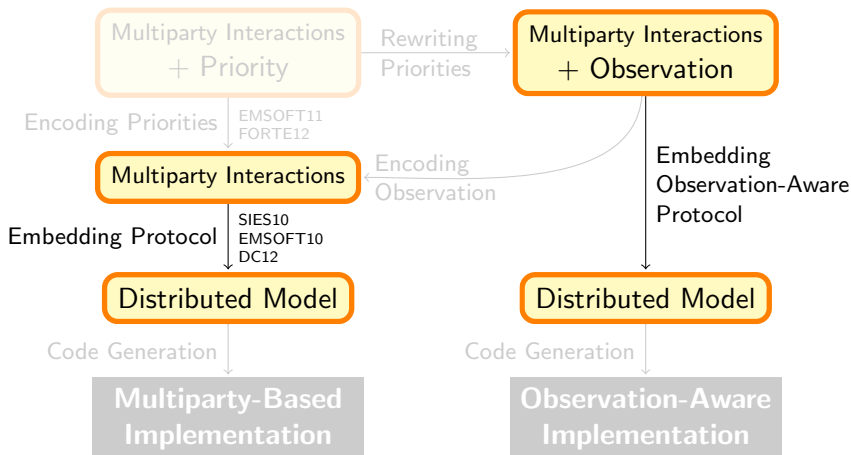


**Multiparty Interactions:**



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## 4 Distributed Protocols



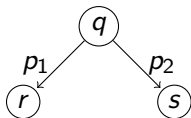
# Distributed components

In a distributed context, atomicity of transitions is broken.

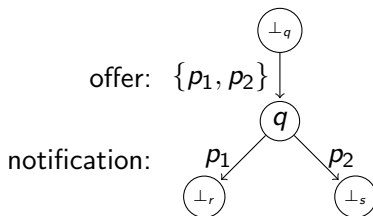
Each transition (in distributed atomic components) is done in two steps:

- sending an offer (indicating enabled ports), through a special port
- receiving a notification (indicating on which port to execute)

**Original Behavior:**



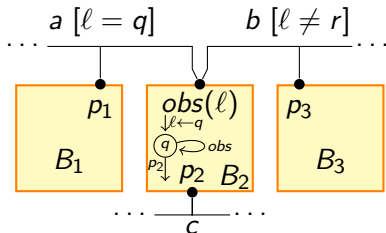
**Distributed Behavior:**



A distributed protocol gathers offer and sends notifications.

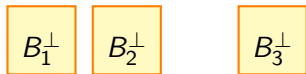
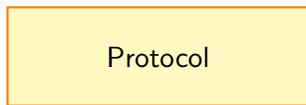
# Distributed Protocol for Multiparty Interactions

## Centralized Example



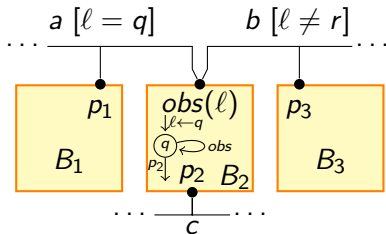
Protocol

## Distributed Version



# Distributed Protocol for Multiparty Interactions

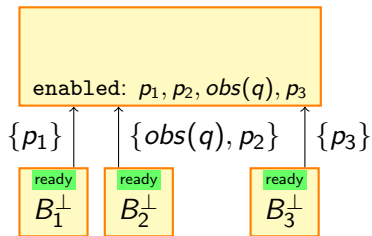
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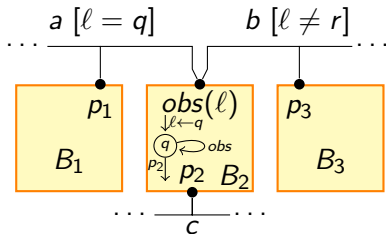
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  - ▶ Mark component as ready
  - ▶ Record enabled ports

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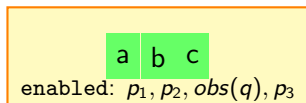


# Distributed Protocol for Multiparty Interactions

## Centralized Example



## Distributed Version

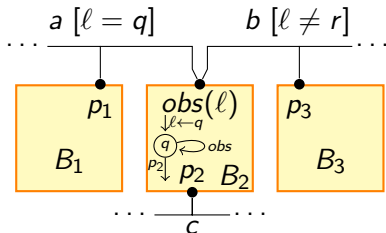


## Protocol

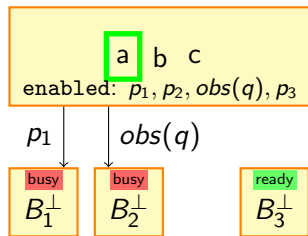
- Gather offers
  - ▶ Mark component as ready
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- Detect enabled interactions
  - ▶ all ports are enabled
  - ▶ all participants are ready

# Distributed Protocol for Multiparty Interactions

## Centralized Example



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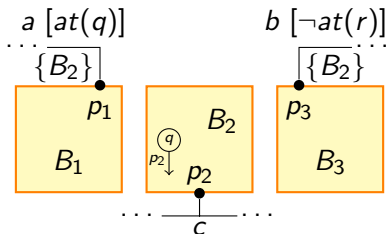


## Protocol

- Gather offers
  - ▶ Mark component as ready
  - ▶ Record enabled ports
- Detect enabled interactions
  - ▶ all ports are enabled
  - ▶ all participants are ready
- Select an enabled interaction and execute it:
  - ▶ Send notifications to participants
  - ▶ Mark participants as not ready

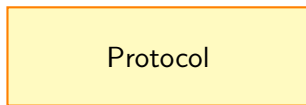
# Observation-Aware Implementation

## Centralized Example



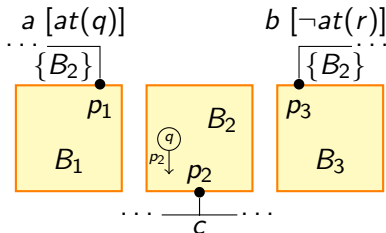
## Protocol

## Distributed Version



# Observation-Aware Implementation

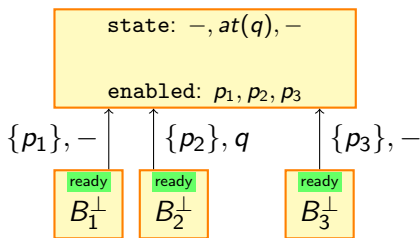
## Centralized Example



## Protocol

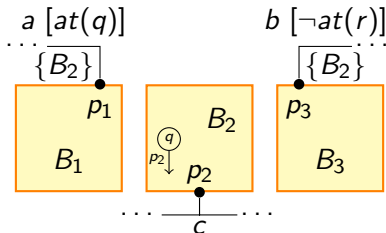
- Gather offers
  - ▶ Mark component as ready
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  - ▶ **Record states**

## Distributed Version

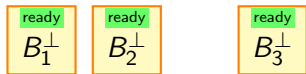
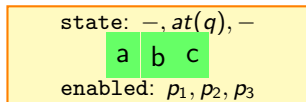


# Observation-Aware Implementation

## Centralized Example



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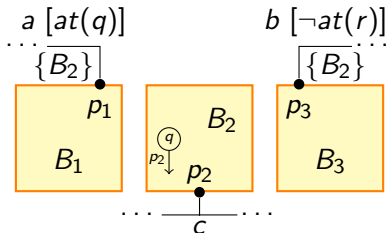


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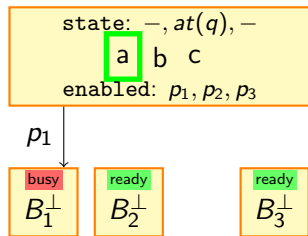
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  - ▶ **Record states**
- Detect possible interactions
  - ▶ all ports are enabled
  - ▶ participants **and observed components** are ready
  - ▶ **state predicate holds**

# Observation-Aware Implementation

## Centralized Example



## Distributed Version

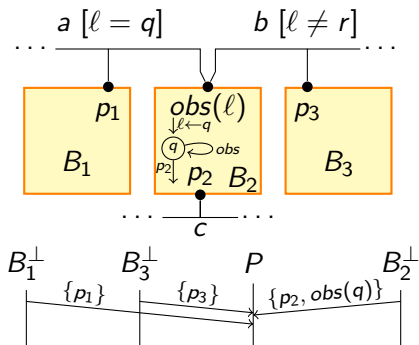


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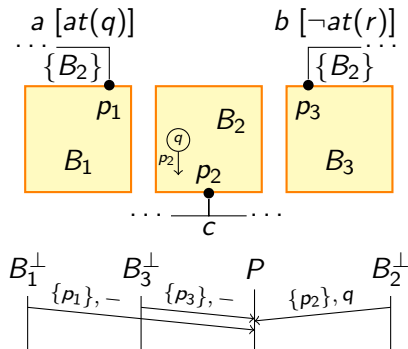
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# Message Efficiency Comparison

## Multiparty-Based

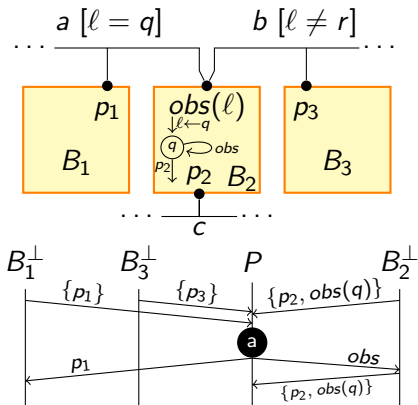


## Observation-Aware

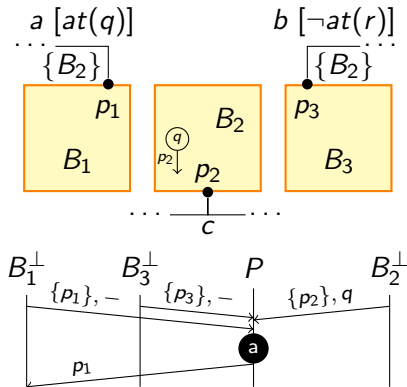


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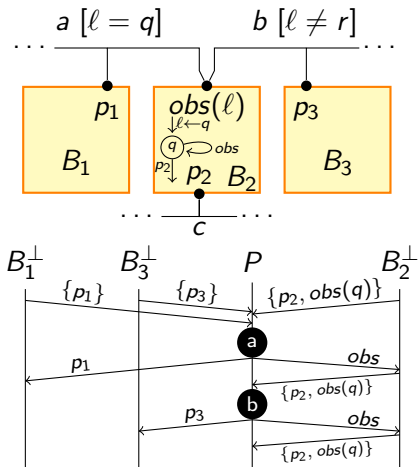


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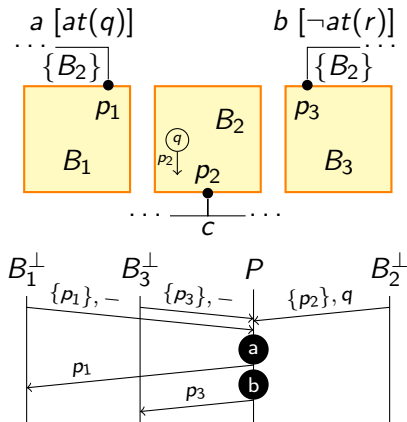


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## Multiparty-Based

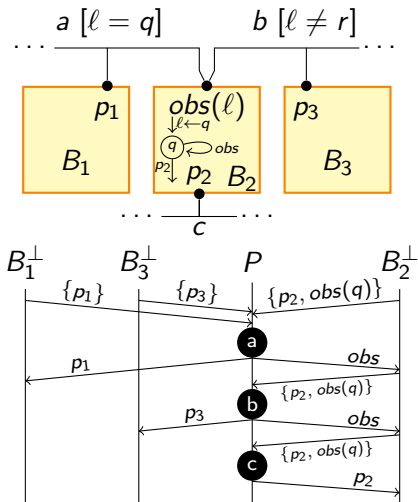


## Observation-Aware

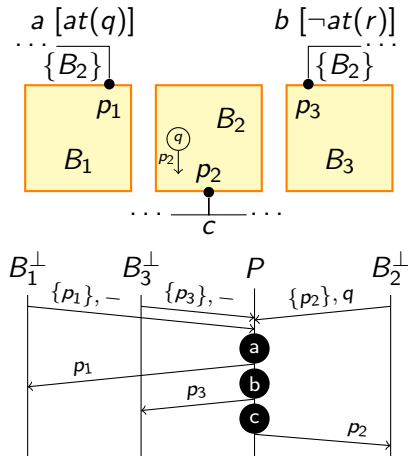


# Message Efficiency Comparison

## Multiparty-Based



## Observation-Aware

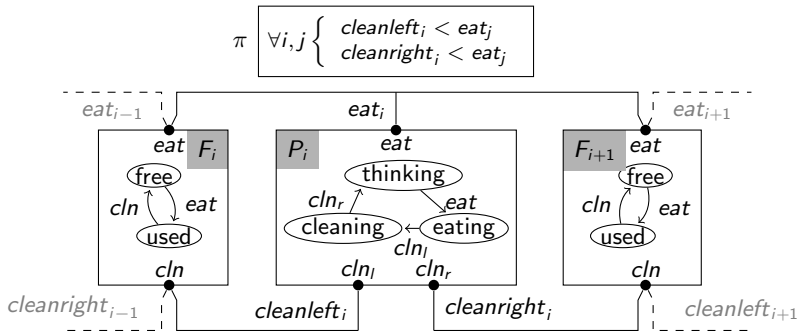


# Outline

- 1 Introduction
- 2 The BIP Model
- 3 Alternative Model: Observation
- 4 Distributed Protocols
- 5 Experiments**
- 6 Conclusion

## Example: Dining philosophers

Model with  $N$  philosophers and  $N$  forks.

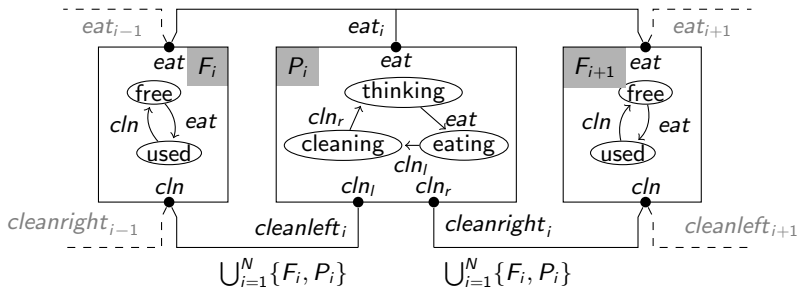


Only interactions *cleanright<sub>i</sub>* and *cleanleft<sub>i+1</sub>* have low priority.

## Example: Dining philosophers

Model with  $N$  philosophers and  $N$  forks.

$$\pi \quad \forall i, j \left\{ \begin{array}{l} \text{cleanleft}_i < \text{eat}_j \\ \text{cleanright}_i < \text{eat}_j \end{array} \right.$$



$$\bigwedge_{j=1}^N \neg(\text{at}(F_j.\text{free}) \wedge \text{at}(P_j.\text{thinking}) \wedge \text{at}(F_{j+1}.\text{free}))$$

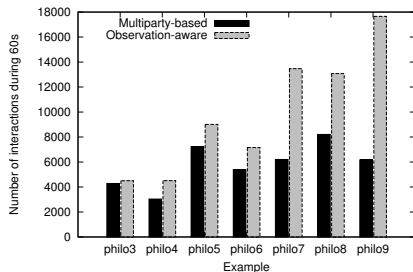
$$\bigwedge_{j=1}^N \neg(\text{at}(F_j.\text{free}) \wedge \text{at}(P_j.\text{thinking}) \wedge \text{at}(F_{j+1}.\text{free}))$$

Only interactions *cleanright<sub>i</sub>* and *cleanleft<sub>i+1</sub>* have low priority.

They need to observe all components.

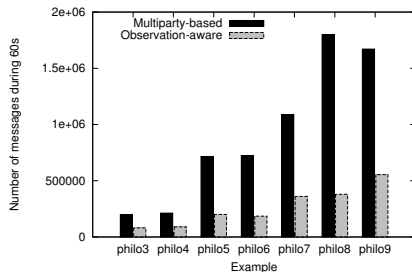
# Dining Philosophers

Interactions executed (during 60s)



Observation-aware implementation gives better performance.

Messages Exchanged (during 60s)

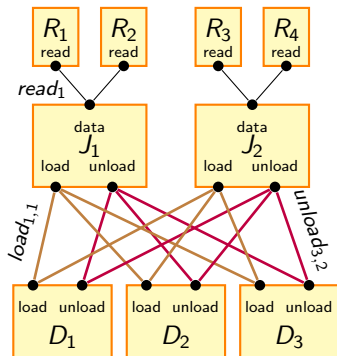


Observation-aware implementation requires fewer messages.

# Jukebox Example

$$\forall i, j, k, \text{unload}_{i,j} < \text{read}_k$$

$$\forall i \in \{2, 3, 4\}, \text{load}_{1,1} < \text{load}_{i,1}$$

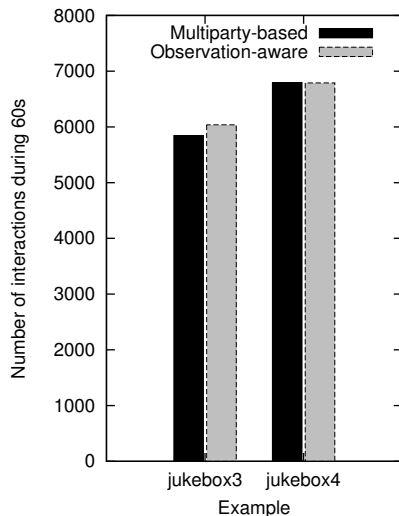


Each reader needs to access disks following a given sequence.

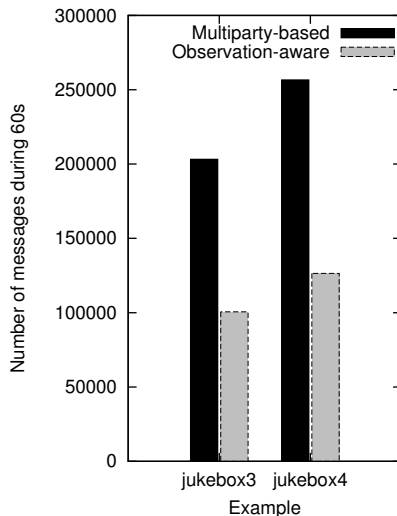
Each jukebox loads and unloads all 3 disks sequentially, in a random order.

An interaction "read" is possible only if the disk loaded by the jukebox is the disk needed by the reader.

Interactions during 60s.



Messages exchanged during 60s.



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## Related Work

### Protocols for Multiparty Interaction:

**Bagrodia** *Process Synchronization: Design and performance evaluation of distributed algorithms.* 1989

**Kumar** *An implementation of n-party synchronization using tokens.* 1990

**Pérez, Corchuelo, Toro** *An order-based algorithm for implementing multiparty synchronization* 2004.

### Distributed implementation for BIP:

[DC12] **Bonakdarpour, Bozga, Jaber, Quilbeuf, Sifakis** *A framework for automated distributed implementation of component-based models.* 2012

[EMSOFT11] **Bonakdarpour, Bozga, Quilbeuf** *Automated distributed implementation of component-based models with priorities.* 2011

### Reducing number of observed components using Knowledge:

[FORTE12] **Bonsalem, Bozga, Quilbeuf, Sifakis** *Knowledge-Based distributed conflict resolution for multiparty interactions and priorities.* 2012

# Conclusion

- We proposed a new glue operator, Observation, for scheduling multiparty interactions.
- We showed that it is expressive enough to encompass Priority.
- We proposed an Observation-Aware protocol that requires fewer messages than a Multiparty-Based protocol for Observation.
- Experiments showed that this message efficiency can improve performance.

## Future Work

- Characterization of the expressiveness allowed by Observation, and comparison with other existing glues.
- Deeper experimental analysis of the new Protocol, in particular when combined with knowledge-based methods:
  - ▶ to reduce number of observed components [FORTE12]
  - ▶ to optimize the protocol itself [HVC'12]

# Thank You